

Principles of Programming Languages

2013.07.24

Notes

- Total available time: 2h.
- You may use any written material you need.
- You cannot use computers or phones during the exam.

1 Prolog (11 points)

1. Define the `prefix` predicate that holds iff its second argument is a prefix of the first argument.
E.g. `prefix("Hello world", "Hello")` is true, while `prefix("Hello world", "wor")` is not.
2. Define an analogous predicate for suffixes.
E.g. `suffix("Hello world", "world")` is true, while `suffix("Hello world", "Hello")` is not.
3. Define the `infix` predicate that holds iff its second argument is a substring of the first argument.
(Hint: an infix is a prefix of a suffix.)
4. Define the `overlap` predicate that holds iff its two argument strings actually overlap, i.e. either one is an infix of the other, or one's prefix is a suffix of the other.

2 Haskell (11 points)

Consider an immutable doubly linked list datatype (`DList`), where each node has two pointers, one to the previous node (`prev`) and another to the next node (`next`), together with its local datum. There is a special value `Nil`, denoting the empty `DList`. A well-formed `DList` has always the first node with `prev` set to `Nil`, and the last node with `next` set to `Nil`.

1. Define the `DList` datatype. `DList` must be an instance of the `Eq` class, and `==` must always terminate for every well-formed `DList`.
2. Define the `car` and `cdr` functions for `DLists`. The latter must return well-formed `DLists`, if not called on `Nil`. Errors must be managed in the `Maybe` monad.
3. Define the `cons` function for `DLists`.

3 Scheme/Ruby (10 points)

Define a *mutable* variant of `DList` either in Scheme or in Ruby. You are requested to define the `DList` datatype; `Dcar`, `Dcdr`, and `Dcons`, i.e. `car`, `cdr`, `cons` variants for `DLists`; and `DList=?` that holds if both its arguments are equal.

Solutions

Prolog

```
prefix([X|_], [X]).
prefix([X|Xs], [X|Z]) :- prefix(Xs, Z).

suffix(X,X) :- \+ X = [].
suffix(_|Xs, S) :- suffix(Xs, S).

infix(X,Y) :- suffix(X, SuffX), prefix(SuffX, Y).

overlaph(X,Y) :- suffix(X, SuffX), prefix(Y, SuffX).

overlap(X,Y) :- overlaph(X, Y); overlaph(Y, X);
                infix(X, Y); infix(Y, X).
```

Haskell

```
data DList a = Nil | Node (DList a) a (DList a)

instance Eq a => Eq (DList a) where
  Nil == Nil = True
  (Node p c n) == (Node p' c' n') = c == c' && n == n'
  _ == _ = False

car Nil = Nothing
car (Node prev head next) = Just head

cdr Nil = Nothing
cdr (Node prev head next) =
  let Node p c n = next
  in Just $ Node Nil c n

cons x Nil = Node Nil x Nil
-- cons exploits call by need: new's definition is recursive:
cons x (Node Nil cur next) = let new = Node Nil x (Node new cur next)
                             in new
```

Scheme

```
(struct DList (prev
              curr
              next) #:mutable)

(define Nil (DList #f #f #f)) ;; the Nil object
(define (Nil? x) (eq? x Nil)) ;; just for convenience

(define (Dcons item node)
  (if (Nil? node)
      (DList Nil item Nil)
      (let* ((newcar (DList Nil (DList-curr node) (DList-next node)))
             (newnode (DList Nil item newcar)))
          (set-DList-prev! newcar newnode)
          newnode)))

(define (Dcar node)
  (if (Nil? node)
      (error "Dcar of Nil")
      (DList-curr node)))

(define (Dcdr node)
  (if (Nil? node)
      (error "Dcdr of Nil")
      (let ((next (DList-next node)))
          (DList Nil (DList-curr next) (DList-next next)))))

(define (DList=? node1 node2)
  (cond
    ((and (Nil? node1)(Nil? node2)) #t)
    ((or (Nil? node1)(Nil? node2)) #f)
    (else
     (and (equal? (DList-curr node1)
                  (DList-curr node2))
          (DList=? (DList-next node1)
                  (DList-next node2))))))
```